DNSSECbis Lookaside Validation

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Topics

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- DNS Delegation and Resolution
- DNSSECbis Data and Traversal
- DLV Overview
- DLV Operations
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Introduction

- DNS in 1987 was to replace HOSTS.TXT and allow for future expansion
- Authenticity of DNS data wasn't considered
- From 1994 to 2005 (and beyond), IETF designs and redesigns Secure DNS
- Secure DNS deployment depends on miracles
- DLV is an early deployment aid

DNS Data and Delegation

• Domain names lay inside a hierarchy of *zones*

- zone boundaries are always at label boundaries
 - but, not all label boundaries are zone boundaries
- every zone except "the root" has ancestors
- any zone can have descendents, by *delegation*
- "root" zone is ultimate ancestor of all zones
- every zone has some *authority* servers
- DNS nodes can contain *resource record sets*
 - sets denoted by <*name,type,class*>
 - each record has some kind of data
 - NS RR set introduces a child zone (*delegation point*)

DNS Traversal and Recursion

• A server is authoritative for zero or more zones

- if zero, server is a "caching forwarder"
- if nonzero, server is an "authority server"
- Authority response types
 - negative: "no name matches your qname"
 - empty: "name is good, but no rrsets match your qtype"
 - positive: "here's what you asked for"
 - referral: "that's in a subdomain, go ask other servers"
- Caching forwarder behaviour
 - acts on behalf of "stub" resolvers
 - caches data for reuse, follows referrals, etc
 - configured to know list of "root" zone servers

DNSSECbis Data and Traversal

• New DNS metadata RR types

- DNSKEY: public key, found at a zone's apex
- RRSIG: generated using RR set data + private key
- NSEC: authenticates unused name space
- DS: in parent zone, authenticates zone's DNSKEY
- Validation
 - Positive answers will include an RRSIG (+ DNSKEY)
 - Referral answers will contain a signed DS (with NS)
 - Negative or empty answers will contain an NSEC
 - Validator must be configured to know trust anchor(s)
 - ultimately this means knowing the public key for "root"

Problems in DNSSECbis Approach

- Trust anchors are *very* widely distributed
 - there's no way to roll out a new key more than once
 - therefore the number of useful anchors is likely "one"
 - and that "one" has to last for the Internet's lifetime

• Root zone stewardship is politically complicated

- signing the root zone requires a strong permanent key
- DNSSECbis depends on trust among root's stewards
- current stewards are not mutually trustful
- Economic benefits of DNSSECbis are unclear
 - adds value for DNS data consumers and producers
 - adds great cost, little revenue for registries/registrars
 - DNS autonomy means "monopoly powers" govern

DLV Overview

- Local policy mechanism for validators
 - not an IETF standard producer/consumer "co-op"
 - only affects results that would have been unsecured
- Early deployment aid
 - supports market growth from 0%, but not full Internet
 - will die after "root" and some large TLDs are secured
- Supports/expects migration to "real DNSSECbis"
 - lets producers/consumers have Secure DNS now
 - creates a market to support registry/registrar costs
 - allows politicos more time to improve stewardship

DLV Metadata

- DLV resource record
 - structurally identical to DS RR (differs semantically)
 - RR type code number is from experimental space
- DLV namespace
 - is within normal DNS namespace
 - normal DNSSECbis is used to secure it
 - can have normal interior zone cuts and delegations
- Example
 - DLV namespace at DLV.ISC.ORG
 - DNSKEY exists for VIX.COM
 - no DS for VIX.COM (or, most likely, for COM)
 - insert DLV RR at VIX.COM.DLV.ISC.ORG

DLV Validation

- Validators are configured with one or more DLV namespaces and trust anchors
- Whenever normal DNSSECbis metadata cannot be found or validated...
 - select the best matching DLV namespace known
 - select the best matching DLV RR within that space
- Examples
 - if two DLV name spaces are known, "root" and MIL
 - no MIL name would ever be searched in the "root" DLV
 - if a DLV namespace knows COM and VIX.COM
 - the VIX.COM DLV would take precedence over COM's for queries of VIX.COM, WWW.VIX.COM, etc

Aggressive Negative Caching

- Possibility of MiM attacks requires that validator issue <u>many</u> DLV queries
- Cached NSEC RRs can obviate these queries
- Problem: NSEC isn't intended for negative caching
- Solution: *Off-The-Wire* negative caching
 - the DLV logic in the validator is "like an application"
 - applications are free to interpret cached NSECs
- Result: most DLV queries will be "suppressed"
- Example
 - cached NSEC declared nonexistence between AAA.DLV.ISC.ORG and CCC.DLV.ISC.ORG
 no need to query for BBB.DLV.ISC.ORG

DLV Operations

- DLV Registry: accept public keys from verified zone owners over repudiable channels; publish
 should be public benefit corporation with cost-based fee structure, who knows when to cease operations
- DLV Registrant: submit to DLV Registry the DNSKEY values from signed zones
 - submissions can cease once the zone's parent is secured, if parent uses DLV or if DLV is dead/dying
- Validator Operators: retrieve and configure trust anchors and DLV namespace info from Registry

 monitor registry in case of key rollover events

Conclusion

- Secure DNS is urgently and much needed
- DNSSECbis works in the lab but isn't deployable
- DLV is an early deployment aid
 - should scale well enough
 - shouldn't scale too well
- DLV is not an IETF standard just a "co-op"
- ISC is committed to DLV
 - will support DLV in BIND9 (9.4.0, due "soon")
 - will operate a robust DLV registry (similar to f-root)
 - will kill DLV when the need for it is past

Questions

• Who else worked on this?

- David Conrad, Johan Ihren, Mark Kosters, Sam Weiler, Mark Andrews, and many others
- Nobody endorses it other than Paul Vixie and ISC
- Why isn't this an IETF protocol?
 - deployment is "just a detail" (ivory-towerism)
- Why did ISC decide to do DLV?
 our mission statement made us do it
- What else do you want to know?