Field Experiments on Post-Quantum DNSSEC

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Steps Taken

- Implemented via liboqs (with regular unassigned algorithm numbers)
 - Falcon512
 - Dilithium2
 - SPHINCS+-SHA256-128s
 - XMSSmt-SHA256-h40-4 / XMSSmt-SHA256-h40-8 (and other parameter sets)
- Measurements using **RIPE ATLAS** (~10,000 probes, ~2M queries in May 2024)
- Deployed BIND9 and Powerdns based zones
- Output variables: Rcode, Correctness, AD bit, response time
- Pre-selection: Exclude ...
 - probe-resolver combinations with incorrect response for RSA-SHA256 (due to noise)
 - resolvers in private IP ranges (due to RIPE ATLAS limitation for TCP)
 - $\circ \quad \text{timeouts and network errors} \\$

Correct responses for a valid label



Percentage of Correct Responses without DO bit

Percentage of Correct Responses with DO bit



Correct responses for a nonexistent label



Queries Using a PQC-aware Resolver

dig +dnssec A dilithium2.pdns.pq-dnssec.dedyn.io @bind9.pq-dnssec.dedyn.io -p 5304

;; Truncated, retrying in TCP mode.

; <<>> DiG 9.18.24-0ubuntu0.22.04.1-Ubuntu <<>> +dnssec A dilithium2.pdns.pq-dnssec.dedyn.io @bind9.pq-dnssec.dedyn.io -p 5304 ;; global options: +cmd

;; Got answer:

;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 22245

;; flags: qr rd ra ad; QUERY: 1, ANSWER: 2, AUTHORITY: 0, ADDITIONAL: 1

;; OPT PSEUDOSECTION:

; EDNS: version: 0, flags: do; udp: 1232

; COOKIE: 8455829f86d7fb7601000000669b5d9517dfc67dff539cac (good)

;; QUESTION SECTION:

;dilithium2.pdns.pq-dnssec.dedyn.io. IN A

;; ANSWER SECTION:

dilithium2.pdns.pq-dnssec.dedyn.io. 3599 IN A 95.217.209.184

;; Query time: 56 msec

;; SERVER: 35.232.14.170#5304(bind9.pq-dnssec.dedyn.io) (TCP)

;; WHEN: Fri Jul 19 23:47:49 PDT 2024

;; MSG SIZE rcvd: 2593

CONTACT 🖸

Post-Quantum DNSSEC Testbed with BIND and PowerDNS

Query our PQC-enabled DNS Resolvers

Send queries to our post-quantum enabled validating resolver! You can choose from a number of post-quantum (and classical) signing schemes, NSEC or NSEC3 mode, and implementations for PowerDNS (source [2]) and BIND (source [2]).

Zones signed accordingly are available at **{algorithm}.{vendor}.pq-dnssec.dedyn.io**, and each has a **A** and a **TXT** record configured. To query a non-existing name, prepend the **nx** label (for example).

Queries will be sent from your browser using DNS-over-HTTPS to a BIND or PowerDNS resolvers with validation support for the selected algorithm. The resolver will talk to the corresponding BIND or PowerDNS authoritative DNS server (again, with support for the selecting signing scheme), to get your response. It will then validate the signature and send the result to your browser.

All queries are send with the DNSSEC_OK flag (+dnssec in dig), so you will see RRSIG and NSEC/NSEC3 records the the responses.



Try it yourself!

https://pq-dnssec.dedyn.io/ (also has detailed results)

What we observed

• Transmission issues are real

- \circ PQC response delivery rates go down significantly as response sizes increase \rightarrow Falcon leads
- Gets worse depending on circumstances, like with DO bit or with NSEC3
- UDP & DO=0:

~70% KSK/ZSK responses correct ~80% CSK responses correct

 \circ $\,$ Goes up by ~10% via TCP $\,$

• UDP & DO=1:

- ~50% responses correct
 - Goes up by ~20-40% via TCP
- 8.5% of probe-resolver pairs claim successfully validating Falcon

The Future? Merkle Trees

- Want to use PQC while keeping messages small
- Use Merkle trees to compress zone signing overhead
 - Signatures become authenticating paths
 - ZSK DNSKEY becomes the root hash
 - $\circ \quad \text{KSK is some secure algorithm with unpleasantly long signatures} \\ \rightarrow \text{replace with small Merkle authenticating path}$

What is a Merkle Tree?

What is a Merkle Tree?

Can we apply this to DNS?

Sure!

- Use a Standardized DNSSEC Algorithm for our KSK
 - Provides Authenticity and Integrity
- Define a new "Merkle Tree" algorithm and store its root hash in the ZSK rdata
 - Provides Integrity via proof of inclusion + gets Authenticity from being signed by KSK
- Signatures become the authenticating path of the Merkle tree
 - Grow logarithmically with the number of RRSets in a zone
- We can combine the work from Batched Signatures Revisited [1] to reduce hash size without reducing security (Second Preimage Resistance)

[1] <u>https://pub.sandboxaq.com/publications/batch-signatures-revisited</u>

We need to change some things about DNS first...

- Circular signing is an issue
 - Everytime you sign something, the Merkle tree changes, and its root node (ZSK) changes
 - Everytime the root node changes, the keytag changes \rightarrow signature's input changes
- We need to allow for our DNSKEY set to have "disjoint" algorithms
 - The DNSKEY set cannot be a part of the Merkle Tree

- We need to change how RRSets are signed when using a Merkle Tree
 - We cannot include the key tag as part of the data being signed
 - We cannot require RRSIGs for all DNSKEY algorithms
 - Instead, DNSKEY is signed with KSK algorithm, and all else signed with ZSK algorithm

We get two nice wins

DNS messages without DNSKEY set stay below line of peril!

Tiny zone transfers

- Since a private key isn't involved, we can have all secondary servers rebuild the tree and authenticating paths
- Interesting trade-off: We can transmit empty signatures during zone transfers greatly reducing the size of the zone
 - Only one signature in zone transfer (for DNSKEY RRset)

We don't have an implementation for this

Was there a difference in the ATLAS tests?

Correct responses for a valid label

Percentage of Correct Responses with DO bit

Correct responses for a nonexistent label

Percentage of Correct Responses without DO bit

Percentage of Correct Responses with DO bit

dnssec-signzone

- Currently only supports offline signing
- Heavy modifications to BIND's dnssec-signzone
 - Iterate through all RRSets and add them to the Merkle tree
 - Finalize the Merkle tree and update keytag
 - Iterate over all RRSIGs and insert the correct authenticating path and keytag
 - Takes about half the time of signing the same zone with ECDSA
 - Additional optimization opportunities might be possible

Some takeaways for Merkle trees

- DNSSEC protocol changes would need to be made
- By defining it with its own algorithm id you can use Merkle trees with any other DNSSEC algorithm
- Zone updates are limited by the root node's (ZSK) TTL
 - Verisign's MTL might help with this?
- DNSKEY messages are not compressed
- Improve deliverability for large signature zones
- Unlike stateful hash based signatures draft no central state is required to be maintained

Thank you!

Acknowledgments:

SSE \uparrow deSEC

Context & Motivation

- In 2022, performed (local-only) DNSSEC study with **Falcon** in PowerDNS
 - Results: <u>https://blog.powerdns.com/2022/04/07/falcon-512-in-powerdns</u>
- Now: Broader experiments with **multiple PQC algorithms**
 - fast validation, short signatures, short-ish keys
- Goal: **Public deployment** on the Internet, to investigate ...
 - **behavior of non-PQC-aware resolvers** typically used by clients
 - behavior of PQC-aware resolvers
- Parameters:
 - KSK/ZSK (BIND) vs. CSK (PowerDNS)
 - Name existence and NSEC vs. [NSEC3 conventional (BIND) vs. minimal (PowerDNS)]
 - UDP vs. TCP
 - DO bit

Algorithm Considerations

Algorithm	NIST Verdict	Approach	Private key	Public key	Signature	Sign/s	Verify/s
Crystals-Dilithium-II [29]	Finalist	Lattice	2.8kB	1.2kB	2.0kB		
Falcon-512 [31]	Finalist	Lattice	57kB	0.9kB	0.7kB	3,307	20,228
Rainbow-I _a [56]	Finalist	Multivariate	101kB	158kB	66B	8,332	11,065
RedGeMSS128 [16]	Candidate	Multivariate	16B	375kB	35B	545	10,365
Sphincs ⁺ -Haraka-128s [11]	Candidate	Hash	64B	32B	8kB		
Picnic-L1-FS [17]	Candidate	Hash	16B	32B	34kB		
Picnic2-L1-FS [17]	Candidate	Hash	16B	32B	14kB		
EdDSA-Ed22519 [12]		Elliptic curve	64B	32B	64B	25,935	7,954
ECDSA-P256 [12]		Elliptic curve	96B	64B	64B	40,509	13,078
RSA-2048 [12]		Prime	2kB	0.3kB	0.3kB	1,485	49,367

Müller, M. et al.: Retrofitting post-quantum cryptography in internet protocols: a case study of DNSSEC. SIGCOMM Comput. Commun. Rev. 50, 49–57 (2020)

- Selected algorithms with public keys and signatures < 10 KB
- Plus: a stateful hash-based algorithm (XMSS)

tcp = False | do = True

0.423%

0.121%

- 10³

- 10²

- 10¹

· 10⁰

- 10³

- 10²

- 10¹

100

99.3%

0.121%

0.123%

tcp = False | do = False

0.185%

99.6%

0.123%

00 unsigned

algo

vendor='pdns', is_nx=True, good-rsa

26

Crypto Algorithm Run Time (PowerDNS)

Backup: Outlook

- Fixing may require **revamping signature representation** in DNS
 - <u>ARRF</u>?
 - Does not necessarily involve a wire format / spec change
 - Or will more robust DoT/DoH/DoQ gain enough traction?
- What would it take to make the root quantum-safe?
 - Further complications from double-signing is this really needed?
- To transition, any scalable solution will require DS provisioning automation
- Future work needed!
 - \rightarrow <u>Research agenda</u>